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CHAPTER XVII: FIXING IMPROPERLY WORKING MODELS

Empirical models do not always yield acceptable solutions. This chapter contains discussion of unacceptable solution conditions and techniques for diagnosing the causes of such conditions.

17.1 Unacceptable Solution Conditions

Four cases of improper solutions can arise. First, a solver could fail exhibiting: a) a time, iteration, or resource limit; b) a lack of meaningful progress; or c) a report of numerical difficulties. Second, a solver may halt identifying that the problem is infeasible. Third, a solver may halt identifying that the problem is unbounded. Fourth, the solver may yield an "optimal," but unacceptable solution.

17.1.1 Solver Failure -- Causes and Prevention

When solvers fail because of numerical difficulties or use an unrealistically large amount of resources to make little progress, the modeler is often in an awkward position. However, several actions may alleviate the situation.

One should first examine whether the model specification is proper. The section on structural checking below gives some techniques for examining model structure. In addition traditional input (commonly called MPS input) based solvers frequently fail because of improper coefficient location (although GAMS prevents some of these errors). In particular, errors can arise in MPS coefficient placement or item naming resulting in more than one (duplicate) coefficient being defined for a single matrix location. Given our concentration on the GAMS modeling system, procedures for finding duplicate coefficients will not be discussed. Nevertheless, this is probably the most common reason why MPS input based solvers run out of time.

The second reason for solver failure involves degeneracy induced cycling. Apparently, even the best solvers can become stuck or iterate excessively in the presence of massive degeneracy. Our experience with such cases indicates one should use an a priori degeneracy resolution scheme as discussed below. We have always observed reduced solution times with this modification.

Thirdly, a solver may fail citing numerical difficulties, an ill-conditioned basis or a lack of progress. Such events can be caused by model specification errors or more commonly poor scaling. Often one needs to rescale the model to narrow the disparity between the magnitudes of the coefficients. Scaling techniques are discussed below.

All of the preventative techniques for avoiding solver failures can be used before solving a model. Modelers should check structure and consider scaling before attempting model solutions. However, degeneracy resolution should not usually be employed until a problem is identified.

17.1.2 Unbounded or Infeasible Solutions

Often the applied modeler finds the solver has stopped, indicating that the model is infeasible or unbounded. This situation, often marks the beginning of a difficult exercise directed toward

finding the cause of the infeasibility or unboundedness, particularly when dealing with large models. There are several techniques one can use when this occurs. The first involves structural checking to find obvious model formulation defects. The second and third techniques involve the use of artificial variables and large upper bounds to find difficulties. Finally one could use the techniques called budgeting and row summing.

17.1.3 Unsatisfactory Optimal Solutions

Unfortunately, optimal solutions can be unrealistic. Discovering an optimal solution means the problem has a mathematically consistent optimum. However, mathematical consistency does not necessarily imply real world consistency (Heady and Candler). Usually, unrealistic solutions may be caused by improper problem specification or assumption violations. Cases arise where the model solution is improper because of: a) omitted constraints or variables; b) errors in coefficient estimation; c) algebraic errors; or d) coefficient placement errors.

Basically, a model may be judged improper because of incorrect valuation or allocation results. Valuation difficulties arise from the reduced cost or shadow price information, such items take on values when primal reduced costs are formed. Allocation difficulties arise when the slack or decision variable values are unrealistic. The values of these items are formed through the constraint interactions. Thus, to diagnose the cause of the unrealistic solution, one investigates either the reduced costs associated with the nonbasic primal variables or the calculations inherent in the primal constraints. Two techniques are presented below, one for the investigation of reduced costs, which we call "budgeting"; and another for the reconstruction of the constraint calculations, which we call "row summing."

17.2 Techniques for Diagnosing Improper Models

Now suppose we turn our attention to the techniques one might use to alleviate model solution difficulties. Table 17.1 presents an array of the possible problems and an indication of the techniques one might use to diagnose such problems.

17.2.1 Simple Structural Checking

There are some simple yet powerful techniques for checking LP formulations, regardless of their presentation method. These fall into two categories: one numerical and one analytical.

17.2.1.1 Analytical Checking

In the case of analytical techniques, consider the problem:

$$\begin{array}{ll}
 \text{Max} & \sum_j c_j X_j \\
 \text{s.t.} & \sum_j a_{ij} X_j \leq b_i \quad \text{for all } i \\
 & \sum_j e_{nj} X_j = d_n \quad \text{for all } n \\
 & \sum_j f_{mj} X_j \geq g_m \quad \text{for all } m \\
 & X_j \geq 0 \quad \text{for all } j
 \end{array}$$

Certain values of these parameters can cause the model to: 1) be infeasible, 2) contain a set of variables that must be zero, 3) contain redundant constraints, and 4) yield an unbounded solution, or 5) contain variables that are always unattractive. Table 17.2 presents a set of cases where model structures will guarantee these properties. Suppose we elaborate on one case which leads to each of the five properties.

A model formulation can cause infeasibility. Suppose in the first constraint, b_i , is less than zero and all the a_{ij} 's in that particular constraint are nonnegative. Obviously this constraint causes the model to be infeasible, since it is impossible for the sum of nonnegative numbers to be less than or equal to a negative number.

Second, it is possible that the constraints require that certain variables be zero. Consider what happens if in the second constraint the right hand side (d_n) equals to zero and all e_{nj} 's are greater than or equal to zero, then every variable with a nonzero coefficient in that constraint must be zero.

There are also cases where the model possesses redundant constraints. Suppose b_i is positive, but all a_{ij} 's are negative or zero; then, clearly, this constraint will be redundant as the sum of negative numbers will always be less than or equal to a positive number.

Checks can also be made for whether the problem is unbounded or contains variables which will never come into the solution. Consider an activity with a positive objective function coefficient which has all nonzero a_{ij} 's negative, all zero e_{nj} 's and all nonzero f_{nj} 's positive. Clearly, then, this variable contributes revenue but relaxes all constraints. This will be unbounded regardless of the numerical values. Further, variables may be specified which will never come into the solution. For example, this is true when c_j is less than 0, all nonzero a_{ij} 's are greater than 0, e_{nj} 's zero, and nonzero f_{nj} 's negative.

These particular structural checks allow one to examine the algebraic formulation or its numerical counterpart. Unfortunately, it is not possible to make simple statements when the constraint coefficients are of mixed sign. In such cases, one will have to resort to numerical checking. All of the procedures above have been automated in GAMSCHCK although they can be programmed in GAMS (See McCarl, 1977).

17.2.1.2 Numerical Model Analysis

Another model analysis methodology involves numerical investigation of the equations and variables. Here, one prints out the equations of a model (in GAMS by using the OPTION LIMROW and LIMCOL command) and mentally fixes variables at certain levels, and then examines the relationship of these variables with other variables by examining the equations. Examples of this are given in the joint products problem above. Numerical model analysis can also be carried out by making sure that units are proper, using the homogeneity of units tests.

Another numerical technique involves use of a "PICTURE" with which coefficient placement and signs can be checked. GAMS does not contain PICTURE facilities, so we do not discuss the topic here, although one is contained in GAMSCHK (see McCarl, 1977).

17.2.2 A Priori Degeneracy Resolution

Degeneracy can cause solvers to cycle endlessly making little or no progress. Solvers like MINOS (Murtaugh and Saunders, 1983) on occasion give messages like "terminating since no progress made in last 1000 iterations" or "Sorry fellows we seem to be stuck." Our experience with such cases indicates one should use an a priori degeneracy resolution scheme adding small numbers to the right hand sides, especially to those constraints which start out with zero or identical right hand sides. The magnitude of the small numbers should be specified so that they are not the same for all rows and so that they do not materially affect the solution. Thus, they might be random or systematically chosen numbers of the order 10^{-3} or 10^{-4} (although they can be larger or smaller depending on the scaling and purpose of the constraints as in McCarl, 1977). We have always observed reduced solution times with this modification. OSL automatically invokes such a procedure.

17.2.3 Altering Units of Constraints and Variables: Scaling

Scaling is done automatically in a number of algorithms including MINOS which is used in GAMS. However, automatic scaling is not always successful. Modelers are virtually always more effective in scaling (Orchard-Hayes). This section explores scaling procedures, discussing the effects on resulting optimal solutions.

Altering the units of constraints and variables improves the numerical accuracy of computer algorithms and can reduce solution time. Scaling is needed when the disparity of matrix coefficient magnitudes is large. An appropriate rule of thumb is, one should scale when the matrix coefficient magnitudes differ in magnitude by more than 10^3 or 10^4 . In other words, scaling is needed if the a_{ij} coefficient with the largest absolute value divided by the coefficient with the smallest nonzero absolute value exceeds 10000. One achieves scaling by altering the formulation so as to convert: a) the objective function to aggregate units (i.e., thousands of dollars rather than dollars), b) constraints to thousands of units rather than units (i.e., one might alter a row from pounds to tons), or c) variables into thousands of units (e.g., transport of tons rather than pounds).

17.2.3.1 Scaling-The Basic Procedure

Given the LP problem

$$\begin{array}{ll} \text{Max} & c_1 X_1 + c_2 X_2 \\ \text{s.t.} & a_{11} X_1 + a_{12} X_2 \leq b_1 \\ & a_{21} X_1 + a_{22} X_2 \leq b_2 \\ & X_1, X_2 \geq 0 \end{array}$$

Suppose one wished to change the units of a variable (for example, from pounds to thousand pounds). The homogeneity of units test requires like denominators in a column. This implies every coefficient under that variable needs to be multiplied by a scaling factor which equals the number of old variable units in the new unit; i.e., if X_j is in old units and X'_j is to be in a new unit, with a_{ij} and a'_{ij} being the associated units.

$$X_j' = X_j / SC_j;$$

where SC_j equals the scaling coefficient giving the new units over the old units

$$a_{ij}' = a_{ij} / (SC_j).$$

The scaling procedure can be demonstrated by multiplying and dividing each entry associated with the variable by the scaling factor. Suppose we scale X_1 using SC_1

$$\begin{aligned} \text{Max} \quad & SC_1 c_1 X_1 / SC_1 + c_2 X_2 \\ \text{s.t.} \quad & SC_1 a_{11} X_1 / SC_1 + a_{12} X_2 \leq b_1 \\ & SC_1 a_{21} X_1 / SC_1 + a_{22} X_2 \leq b_2 \\ & X_1 \qquad \qquad X_2 \geq 0 \end{aligned}$$

or substituting a new variable $X_1' = X_1 / SC_1$ we get

$$\begin{aligned} \text{Max} \quad & SC_1 c_1 X_1' + c_2 X_2 \\ \text{s.t.} \quad & SC_1 a_{11} X_1' + a_{12} X_2 \leq b_1 \\ & SC_1 a_{21} X_1' + a_{22} X_2 \leq b_2 \\ & X_1' \qquad \qquad X_2 \geq 0 \end{aligned}$$

Variable scaling alters the magnitude of the solution values for the variables and their reduced cost as we will prove later.

Scaling can also be done on the constraints. When scaling constraints; e.g., transforming their units from hours to thousands of hours, every constraint coefficient is divided by the scaling factor (SR)

as follows:

$$\begin{aligned} \text{Max} \quad & c_1 X_1 + c_2 X_2 \\ & a_{11} / SR X_1 + a_{12} / SR X_2 \leq b_1 / SR \\ & a_{21} X_1 + a_{22} X_2 \leq b_2 \\ & X_1, \qquad X_2 \geq 0 \end{aligned}$$

where SR is the number of old units in a new unit and must be positive. Constraint scaling affects :1) the slack variable solution value, which is divided by the scaling factor; 2) the reduced cost for that slack, which is multiplied by the scaling factor; and 3) the shadow price, which is multiplied by the scaling factor.

The way scaling factors are utilized may be motivated by reference to the homogeneity of units section. The coefficients associated with any variable are homogeneous in terms of their denominator units. Thus, when a variable is scaled, one multiplies all coefficients by a scaling factor (the old unit over the new unit) changing the denominator of the associated coefficients. Constraints, however, possess homogeneity of numerator units so, in scaling, we divide through by

the new unit divided by the old unit. Thus, when changing a constraint from pounds to tons one divides through by 2000 (lbs/tons).

Two other types of scaling are also relevant in LP problems. Suppose that the right hand sides are scaled, i.e., scaled from single units of resources available to thousands of units of resources available. Then one would modify the model as follows:

$$\begin{array}{ll} \text{Max} & c_1 X_1 + c_2 X_2 \\ \text{s.t.} & a_{11} X_1 + a_{12} X_2 \leq b_1 / SH \\ & a_{21} X_1 + a_{22} X_2 \leq b_2 / SH \\ & X_1, X_2 \geq 0 \end{array}$$

The net effects of this alteration will be that the optimal value of every decision variable and slack would be divided by the scaling factor, as would the optimal objective function value. The shadow prices and reduced costs would be unchanged.

One may also scale the objective function coefficients by dividing every objective function coefficient through by a uniform constant (SO).

$$\begin{array}{ll} \text{Max} & c_1 / SO X_1 + c_2 / SO X_2 \\ & a_{11} X_1 + a_{12} X_2 \leq b_1 \\ & a_{21} X_1 + a_{22} X_2 \leq b_2 \\ & X_1, X_2 \geq 0 \end{array}$$

Under these circumstances, the optimal decision variables and slack solutions will be unchanged; but both the shadow prices and reduced costs will be divided by the objective function scaling factor as will be the optimal objective function value.

Scaling may be done in GAMS. Namely putting in the statement `variablename.scale = 1000` would cause all variables in the named variable block to be scaled by 1000 with the solution automatically being readjusted. Similarly `equationname.scale = 1000` will scale all constraints in a block. This must be coupled with the command `Modelname.scaleopt=1`.

17.2.3.2 Mathematical Investigation of Scaling

In this section an investigation will be carried out on the effects of scaling using the matrix algebra optimality conditions for a linear program. Readers not interested in such rigor may wish to skip to the summary and empirical example.

The optimality conditions for the LP problem are given by

$$\begin{array}{l} C_B B^{-1} a_j - c_j \leq 0 \text{ for all } j \\ B^{-1} b \leq 0. \end{array}$$

Given such a solution the optimal decision variables are given by

$$X_B = B^{-1} b,$$

the shadow prices by

$$U = C_B B^{-1}$$

and the reduced costs by

$$C_B B^{-1} a_j - c_j,$$

and the optimal Z value is

$$Z = C_B B^{-1} b$$

In our investigation, we examine the impact of scaling on each of these items.

17.2.3.2.1 Variable Scaling

When a variable is scaled, the problem becomes:

$$\begin{array}{ll} \text{Max} & SC_1 c_1 X_1' + c_2 X_2 \\ \text{s.t.} & SC_1 a_{11} X_1' + a_{12} X_2 \leq b_1 \\ & SC_1 a_{21} X_1' + a_{22} X_2 \leq b_2 \\ & X_1 \quad X_2 \geq 0 \end{array}$$

where X_1 equals SCX_1' and SC is a positive scalar.

The effect on the solution depends on whether the scaled variable is basic or nonbasic. First, consider nonbasic variables. If a nonbasic variable is scaled, then the scaling operation does not affect the basis inverse. Thus, the only thing that needs to be investigated is whether or not scaling the nonbasic variable renders it attractive to bring into the basis. This involves an investigation of the reduced cost after scaling. Constructing the reduced cost for this particular variable

$$C_B B^{-1} SC a_j - SC c_j = SC (C_B B^{-1} a_j - c_j)$$

we find that the reduced cost after scaling (new) equals the reduced cost before scaling (old) times the scaling factor. Thus, we have the old reduced cost multiplied by the scaling constant and under positive SC the before scaling solution remains optimal. The only alteration introduced by scaling a nonbasic variable is that its reduced cost is multiplied by the scaling factor. This can be motivated practically. If it costs \$50 to enter one acre of a crop not being grown into solution, it would logically cost \$50,000 to bring in a thousand acres of that crop.

Now suppose a basic variable is scaled. In this case, the basis inverse is altered. Suppose that the basis matrix before scaling is B , while the matrix of technical coefficients before scaling is A . The new matrices (B^*, A^*) can be expressed as the old matrices (B, A) post-multiplied by matrices K_A and K_B which are modified identity matrices. Assuming the n^{th} column of B is being scaled, then the element on the diagonal in the n^{th} column of the K_B matrix will be the scaling factor. Thus,

$$A^* = AK_A \quad B^* = BK_B$$

where

$$K_B = \begin{bmatrix} 1 & 0 & 0 & \dots & 0 & \dots & 0 \\ 0 & 1 & 0 & \dots & 0 & \dots & 0 \\ 0 & 0 & 1 & \dots & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & SC & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & 0 & \dots & 1 \end{bmatrix}$$

The K_A matrix would be formed similarly with the column in the A matrix being scaled identifying the diagonal element where SC appears.

We may derive a relationship between the basis inverses before and after scaling. Matrix algebra theory shows that

$$(B^*)^{-1} = (BK_B)^{-1} = K_B^{-1} B^{-1}$$

We should also note that the scaled objective function coefficients of the basic variables are post-multiplied by K_B , i.e.,

$$C_B = C_B K_B$$

Now let us look at the optimality criteria for non-basic variables

$$C_B^* B^{*-1} a_j^* - c_j^*$$

The reduced cost after scaling becomes

$$\begin{aligned} C_B^* B^{*-1} a_j^* - c_j^* &= C_B K_B K_B^{-1} a_j - c_j \\ &= C_B B^{-1} a_j - c_j \geq 0 \end{aligned}$$

since $K_B K_B^{-1} = I$. Thus, the reduced costs after scaling equal the reduced costs before scaling. Thus, we have proven that the solution will remain optimal.

We now need to turn our attention to whether or not the basic variables remain nonnegative. The values of the basic variables at optimality are given by $B^{*-1}b$. Substituting in our relationships, we obtain

$$X_B^* = B^{*-1} b = K_B^{-1} B^{-1} b = K_B^{-1} X_B.$$

The scaled solution equals K_B^{-1} times the unscaled solution. The inverse of K_B is an identity-like matrix with one over the scaling factor on the diagonal in the position where the scaled variable

enters the basis.

$$K_B^{-1} = \begin{bmatrix} 1 & 0 & 0 & \dots & 0 & \dots & 0 \\ 0 & 1 & 0 & \dots & 0 & \dots & 0 \\ 0 & 0 & 1 & \dots & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & 1/SC & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & 0 & \dots & 1 \end{bmatrix}$$

The $K_B^{-1} X_B$ multiplication yields the vector

$$X_B = K_B^{-1} X_B = \begin{bmatrix} X_{B_1} \\ X_{B_2} \\ \dots \\ X_{B_K} / SC \\ \dots \\ X_{B_m} \end{bmatrix}$$

Thus, scaling reduces the magnitude of the particular basic variable being scaled, while all other variables are unaffected.

We may also investigate the objective function consequences. The optimal objective function value of the scaled problem is

$$Z^* = C_B^* B^{*-1} b = C_B^* K_B K_B^{-1} B^{*-1} b = C_B^* B^{-1} b = Z$$

Clearly, then, the objective function value after the change equals the objective function value before the change. All in all, column scaling leaves the problem with the same qualitative answer. The solution value of the particular variable being scaled and its reduced cost are altered by the scaling factor.

17.2.3.2.2 Effects of Constraint Scaling

When one scales a constraint, the resultant problem appears as

$$\begin{array}{rcll} \text{Max} & c_1 X_1 & + & c_2 X_2 \\ & a_{11} / SR X_1 & + & a_{12} / SR X_2 \leq b_1 / SR \\ & a_{21} X_1 & + & a_{22} X_2 \leq b_2 \\ & X_1 & , & X_2 \geq 0 \end{array}$$

or

$$\begin{array}{ll} \text{Max} & CX \\ \text{s.t.} & RAX \leq Rb \\ & X \geq 0 \end{array} \quad \text{or} \quad \begin{array}{ll} \text{Max} & CX \\ \text{s.t.} & A^*X \leq b^*, \\ & X \geq 0 \end{array}$$

where R is a scaling matrix of the form

$$R = \begin{bmatrix} 1 & 0 & 0 & \dots & 0 & \dots & 0 \\ 0 & 1 & 0 & \dots & 0 & \dots & 0 \\ 0 & 0 & 1 & \dots & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & 1/RS & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & 0 & \dots & 1 \end{bmatrix}$$

Further, the new basis (B^*) is related to the old basis as follows

$$B^* = RB$$

and the basis inverse is the old basis inverse multiplied by the inverse of R.

$$B^{*-1} = B^{-1} R^{-1}$$

Again, R^{-1} is an identity-like matrix quality.

$$R^{-1} = \begin{bmatrix} 1 & 0 & 0 & \dots & 0 & \dots & 0 \\ 0 & 1 & 0 & \dots & 0 & \dots & 0 \\ 0 & 0 & 1 & \dots & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & RS & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & 0 & \dots & 1 \end{bmatrix}$$

Now let us turn our attention to the effects of scaling a constraint. We will derive the results assuming the slack variable is not in the basis. The reduced cost criteria for the scaled problem is given by

$$C_B^* B^{*-1} a_j^* - c_j^* = C_B B^{-1} R^{-1} R a_j - c_j = C_B B^{-1} a_j - c_j$$

Thus, the optimality conditions are the same as before scaling and the solution remains optimal with unchanged reduced costs. We now may investigate the variable solution values. For the

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scaled problem, the solution values are

$$X_B^* = B^{-1} b^* = B^{-1} R^{-1} R b = B^{-1} b$$

which shows that the values of the basic variables are unaffected. The objective function is also unchanged.

Thus, the optimality and feasibility of the basis remain entirely unaffected. What then does change? The shadow prices after scaling are

$$U = C_B B^{-1} R^{-1}.$$

Given the form of R^{-1} from above, the shadow prices are identical to the shadow prices before scaling

$$U_i = (C_B B^{-1})_i$$

for all rows but the particular row being scaled. For that row, the shadow price is multiplied by the scaling factor

$$U_i = (C_B B^{-1})_i RS$$

Finally we should note that when a slack variable is in the basis, then constraint scaling simply changes the magnitude of the optimal slack variable value by dividing it by the scaling factor.

17.2.3.2.3 Objective Function and Right Hand Side Scaling

If the objective function coefficients are uniformly divided by a constant, the values of the solution variables ($B^{-1} b$) are unaffected. However, the magnitudes of the shadow prices ($C_B B^{-1}$), the optimal objective function value ($C_B B^{-1} b$), and the reduced costs ($C_B B^{-1} a_j - c_j$) are affected. In all of these cases, these items would be uniformly divided by the objective function scaling factor. A similar observation can be made regarding scaling the right hand side. The right hand side determines only the objective function value ($C_B B^{-1} b$) and the solution value of the basic variables ($B^{-1} b$). Dividing all right hand sides by a constant would divide the objective function and all the optimal variable values by the same constant.

17.2.3.4 Summary

Scaling alters coefficient magnitudes within the matrix and the resultant magnitude of selected items in the solution. Consider the LP problem

$$\begin{array}{ll} \text{Max} & CX \\ \text{s.t.} & AX \leq b \\ & X \geq 0 \end{array}$$

Suppose a set of positive (all scaling factors must be positive) scaling factors will be applied to the LP model. The scaling factors are a) COLSCAL_j for the jth variable - a factor multiplying every coefficient under that variable, b) ROWSCAL_i for the ith constraint - a factor dividing every coefficient in that constraint, c) OBJSCAL for the objective function - a factor dividing every coefficient in the objective row and d) RHSSCAL for the right hand side - a factor dividing every right hand side value. The parameters of the model after scaling are:

$$\begin{aligned}
 c'_j &= c_j * \frac{\text{COLSAL}_j}{\text{OBJSCAL}} \\
 a'_{ij} &= a_{ij} * \frac{\text{COLSAL}_j}{\text{ROWSCAL}_i} \\
 b'_i &= b_i * \frac{1}{\text{ROWSCAL}_i * \text{RHSSCAL}}
 \end{aligned}$$

where the ' denotes the new coefficients. The relationship between solution items before and after scaling is given in Table 17.3. Thus, if a particular variable j is scaled by 1000 and constraint i is scaled by a 1000, the a_{ij} value is numerically unchanged. However, if variable j is scaled by 1000 and constraint i was scaled by 100, then the value of the a_{ij} coefficient is multiplied by 10.

Thus, for example, if an optimal variable value before scaling was 6 and the right hand side is multiplied by 100 while the coefficients of that variable are multiplied by .02 then the resultant value after scaling would be 3.

Finally, we must emphasize that the only proper way of scaling is to operate on all coefficients for each variable, right hand side and objective function in the same manner. One cannot selectively scale selected coefficients. The interaction of the various scaling factors can make it look like one is only scaling selected coefficients, as will be demonstrated in the empirical example below. But this is not the case, consistency must be maintained.

17.2.3.5 Empirical Example of Scaling

The previous discussion deals with the implications of changes in constraint, variable, right hand side and objective function units. The reader, however, should note that gains in terms of numerical stability may arise only when several items are simultaneously scaled. This is illustrated in the example. Consider the following problem

$$\begin{array}{rllll}
 \text{Max} & X_1 - & 500X_2 & - & 400X_3 & - & 500X_4 & & & \\
 \text{s.t.} & X_1 - & 10000X_2 & - & 8000X_3 & & & \leq & 0 & \\
 & & 5X_2 & + & 4X_3 & - & 50X_4 & \leq & 0 & \\
 & & 1500X_2 & + & 2000X_3 & & & \leq & 600000 & \\
 & & 50X_2 & + & 45X_3 & & & \leq & 30000 & \\
 & X_1 & X_2 & & X_3 & & X_4 & \geq & 0 &
 \end{array}$$

where:

- X₁ is the sale of nibbles in pounds. It returns \$1 of profit per unit and removes one pound from the nibble balance row.
- X₂ is the hours of nibble production via process 1. One hour's worth of production uses \$500 worth of direct cost, 5 units of gribbles, 1500 hibble and 50 hours of labor. As a result, one gets 10,000 nibbles.
- X₃ is the hours of nibble production using process 2. Here 8000 nibbles are produced at a direct cost of \$400 with four pounds of gribbles used, 2,000 hibble, and 45 hours of labor produced.

negative objective function coefficient (when the objective is to maximize) and positive in a single constraint. Artificial variables only remain in the solution when the restrictions with which they are associated cannot be met, as occurs in misspecified LP models. For example, a model might contain a minimum requirement of 10,000 units production whereas the labor resource availability constraint permits fewer units. This problem may arise if: a) the 10,000 unit requirement is too large and has been improperly entered, b) the labor endowment is erroneously too small, c) the labor requirements for production have been overestimated, or d) the contribution to the minimum requirement constraint is too small.

The last three cases arise when the minimum requirement level is correct, but the infeasibility is caused by misspecifications in other coefficients. Thus, infeasibilities arise not only because of improperly specified minimum requirement rows, but also because of errors in other coefficients.

The question is "how can one discover the cause of the infeasibility?" This can be done by adding artificial variables to the model formulation. The inclusion of artificial variables permits all models to have feasible optimal solutions regardless of whether the "real" constraints are satisfied. Infeasible solutions exhibit nonzero artificial variables. Nonzero artificial variables will cause a large negative objective function value and large shadow prices since some C_B 's in the $C_B B^{-1}$ computations are large. Specifically, constraints which exhibit large shadow prices are those involved with the infeasibility. The constraints not causing the infeasibility will have unaffected shadow prices. Thus, the imposition of artificial variables allows one to identify which constraints are nominally causing the infeasibility. We do not argue that such information cannot be found in an ordinary infeasible solution; however, it is more difficult to interpret. Ordinarily, infeasible solver solutions are detected by phase 1 of the simplex algorithm wherein the shadow prices give the marginal contribution of a change in the right hand side to the sum of the infeasibilities.

To illustrate the use of artificial variables in the context of an infeasible model consider the following example:

$$\begin{array}{rll}
 \text{Max} & 50X_1 & + \quad 50X_2 \\
 \text{s.t.} & X_1 & + \quad X_2 \leq 50 \\
 & 50X_1 & + \quad X_2 \leq 65 \\
 & X_1 & \geq 20 \\
 & X_1 & , \quad X_2 \geq 0
 \end{array}$$

This problem is infeasible due to the interaction of the second and third constraints. Suppose that an error was made and the number 50 which is specified as the requirement of X_1 for the second resource should have been 0.50. The third constraint has a minimum requirement, thus an artificial variable is included.

$$\begin{array}{rll}
 \text{Max} & 50X_1 & + \quad 50X_2 & - \quad 10000A \\
 & X_1 & + \quad X_2 & \leq 50 \\
 & 50X_1 & + \quad X_2 & \leq 65 \\
 & X_1 & & A \geq 20
 \end{array}$$

Here the artificial variable A is entered with a large negative number in the objective function and a plus one in the third constraint, thus permitting a minimum requirement to be satisfied. The

solution to the augmented problem is shown in Table 17.8. The value of the objective function is -186,935. In this solution, the artificial variable A is nonzero with the second and third constraints binding. The shadow prices on the second and third constraints reflect the influence of the artificial variable. Thus, the modeler would receive signals that there was something wrong in the interaction of the second and third constraints. Hopefully then the data error would be found.

In summary, artificial variables are useful in finding the source of infeasibility. Artificials are only needed in constraints that are not satisfied when the decision variables are zero. Their use allows the model user to find infeasibilities by narrowing attention to the constraints which are the underlying causes of the infeasible solution. We also feel they should be used where infeasibilities can arise in models which: a) continually have their data altered; and b) are used by people other than the modeler (see McCarl et al.).

17.2.5 Use Unrealistically Large Upper Bounds to Find Causes of Unboundedness

LP problems may also yield unbounded solutions. One can again use structural checking or a picture to find the problem. However, if these checks fail, imposition of large upper bounds in an unbounded model on all variables which exhibit desirable objective function coefficients will prevent unboundedness, and will cause the variables causing unboundedness to take on large solution values. Investigation of the variables which take on such large values will allow the modeler to find the cause of the unboundedness. Consider a LP problem which has large upper bounds imposed.

$$\begin{array}{ll} \text{Max} & cX \\ \text{s.t.} & AX \leq b \\ & X \leq M \\ & X \geq 0 \end{array}$$

Here the equations $X \leq M$ are upper bound constraints limiting the decision variables to a large number (e.g., the constraint $X_1 \leq 100,000$ has been imposed on the model). Given decision variables, M would be set so it was unrealistically large (i.e., 1,000 times larger than the largest expected X value). Why would anyone want to impose such bounds? Consider the following simple example.

$$\begin{array}{llll} \text{Max} & 3X_1 & - & X_2 & + & X_3 \\ & X_1 & - & X_2 & & = & 0 \\ & & & & & & X_3 \leq 5 \\ & X_1, & & X_2, & & & X_3 \geq 0 \end{array}$$

This problem is unbounded: the model can purchase X_2 , using it to produce X_1 at a net operating profit of \$2 without limit. However, the imposition of the constraint $X_1 \leq 100,000$ yields the solution $X_1 = X_2 = 100,000$, $X_3 = 50$. Thus, if the model user saw this solution and felt that $X_1 = 100,000$ was unrealistically large then this would show that there is something wrong within the model. It also shows that X_1 and X_2 are the items involved with the unboundedness while X_3 is not a factor.

The use of large upper bounds precludes the possibility of an unbounded solution but causes the objective function and some of the variables to take on unrealistically large values. Subsequently,

one can trace the cause of the unboundedness by examining the set of variables which are unrealistically large. This is important since ordinarily LP solvers are implemented so that when they discover an unbounded solution they automatically stop. This leaves the user without much information as to the cause of the unboundedness.

Decision modelers may wish to upper bound the sum of a number of variables rather than each and every variable. This could be done by using the following:

$$\begin{aligned} \text{Max} \quad & \sum_j c_j X_j \\ \text{s.t.} \quad & \sum_j a_{ij} X_j \leq b_i \quad \text{for all } i \\ & \sum_j X_j \leq M \\ & X_j \geq 0 \quad \text{for all } j \end{aligned}$$

Here one specifies that the sum of all variables is less than or equal to an unrealistically large number.

17.2.6 Budgeting

Yet another model analysis technique, particularly when dealing with unrealistic optimal solutions, involves Budgeting. Budgeting herein refers to the reconstruction and examination of reduced cost and shadow price information. The procedure is best illustrated through example. Consider the model shown in Table 17.9.

This model contains activities for buying miscellaneous inputs; selling corn, soybeans, and pork; and producing corn, soybeans, and hogs. The model is maximized subject to a resource constraint on land, along with supply-demand balances on pork, soybeans, corn, and miscellaneous inputs. The miscellaneous input item is specified in dollars and therefore enters the objective function at a per unit cost of \$1 while supplying a dollar's worth of miscellaneous inputs. Corn is sold for \$2.50 per unit, soybeans \$6 per unit, and pork \$.50 per unit. Corn production incurs \$75 in direct production costs and \$125 in miscellaneous inputs while using one acre of land and yielding 120 bushels of corn. Soybean production costs \$50 in direct production costs and another \$50 in miscellaneous inputs while using an acre of land and yielding 50 bushels of soybeans. Hog production has no direct costs, uses \$20 in miscellaneous inputs, and requires 20 bushels of corn. An unrealistically large yield in the hog activity has been entered (1000 pounds per hog). This example "error" will be sought by the budgeting technique.

The optimum solution to this model is shown in Table 17.10. The optimal value of the objective function is \$1,508,000. This solution includes several symptoms that there is something wrong. For example, 3,600,000 pounds of pork are sold, the reduced cost on raising soybeans is \$2,480 an acre, the shadow price on land is \$2,680 and the shadow price of corn is \$24 a bushel. Budgeting investigates the shadow prices and reduced costs in an effort to discover model misspecifications.

The matrix and summation formula for reduced costs is

$$C_B B^{*-1} a_j - c_j = \sum_i (C_B B^{*-1})_i a_{ij} - c_j \geq 0,$$

which will be greater than or equal to zero for all nonbasic variables and exactly equal to zero for basic variables. Utilizing the fact that the $C_B B^{-1}$ are the shadow prices, then the equation becomes

$$\sum_i U_i a_{ij} - c_j \geq 0$$

where U_i is the shadow price associated with row i , the a_{ij} 's are the technical coefficients of the original model, and the c_j is the original objective function coefficients associated with X_j .

Budgeting involves construction of an extensive version of the reduced cost calculations which in turn are examined for plausibility. The variable budgeted first could be chosen because: a) it is nonbasic when intuition suggests it should be basic; b) it has an unrealistically high reduced cost; or c) it uses a resource which appears to be improperly valued. In the example, suppose we budget soybean production because of its high reduced cost. To budget, write a row for each nonzero coefficient (a_{ij}) under the chosen variable, with associated shadow prices (U_i) and $a_{ij}U_i$ the product, then finally subtracting cost. The budget for soybean production is shown in Table 17.11.

Mechanically the budget examines the cost of resource usage in those rows for which the activity uses resources and values (shadow price) of these resources. In turn the objective function value of the variable is considered and the reduced costs reconstructed. The soybean production variable has non-zero coefficients in the land, soybean production, and miscellaneous input constraints. The shadow price for land is \$2,680. Thus, one acre of soybeans uses \$2,680 worth of land and yields 50 bushels, each selling for \$6. Also, 50 units of miscellaneous inputs are used which, when valued at \$1, cost \$50. Summing these terms, the marginal contribution of soybean production, ignoring its direct costs, is \$2,430. Its direct cost ($c_j=50$) is then subtracted yielding a \$2,480 reduced cost. One may conclude that, the \$2,480 reduced cost is caused by the \$2,680 shadow price on land. The question then becomes why is land this valuable.

Shadow prices are derived from setting the reduced costs of basic variables to zero. Thus, the high land shadow price must arise from the reduced costs of some basic variable which utilizes land. The only land using basic variable is corn production. We then budget the corn production variable (Table 17.12). Note that while one acre of corn production uses \$2,680 of land, it receives \$2,880 from the value of the corn sold. Here, the reason for the \$2,680 cost of land is the \$2,880 value of the corn. Institutional knowledge indicates the 120 bushels per acre corn yield is reasonable, but the \$24 corn shadow price per bushel is not. Thus, the question becomes, "Why is the corn shadow price so high?" Again, this will be determined by a basic variable which utilizes corn. The only basic cornusing variable is hog production. The budget for hog production is shown in Table 17.13. These computations show that zero reduced cost for this activity requires that 20 bushels of corn be valued at \$24/unit. The cause of the \$500/bushel value for corn is an unrealistic value of pork produced (\$500). The erroneous 1000 pound coefficient for pork production per hog would then be discovered. A revised value of the pork yield per hog would alter the model, making the solution more realistic.

The budgeting technique is useful in a number of settings. Through its use, one may discover why variables are nonbasic when they should be basic. The soybean production variable budget provides such an example. Budgeting, in such a case, may discover difficulties in the particular variable being budgeted or in shadow prices.

Budgeting may also be used to discover why particular activities are basic when modeler intuition suggests they should be nonbasic. For example, by tracing out the costs and returns to corn as opposed to soybean production to see what the major differences that lead to corn being profitable while soybeans are not.

The third use of budgeting involves discovering the causes of improper shadow prices. Shadow prices arise from a residual accounting framework where, after the fixed revenues and costs are considered, the residual income is attributed to the unpriced resources.

Budgeting can also be used to deal with infeasible solutions from Phase I of a Phase I/Phase II simplex algorithm. Phase I of such algorithms minimizes the sum of infeasibilities. Thus, all of the objective function coefficients of the decision variables in the model are set to zero. The phase I shadow prices refer to the amount by which the sum of the infeasibilities will be reduced by a change in the right hand sides. Budgeting then can be done to trace shadow price origins and to see why certain variables do not come into solution. Solutions containing artificial variables may also be budgeted.

17.2.7 Row Summing

Model solutions also may be analyzed by examination of the primal allocation results. In the budgeting example problem, one could have examined the reasons for the sale of 3.6 million pounds of pork. This can be done through a procedure we call row summing. This is illustrated through a slightly different, but related, example Table 17.14.

Compared to the model shown in Table 17.9, the pork production coefficient has been altered to -150, while the corn yield per unit has been changed to an incorrect value of -1200 -- the error. We have also introduced a RHS of 20 on the corn balance equation. The solution to this model is shown in Table 17.15. The optimal value of the objective function is \$1,860,055. Here 5.4 million pounds of pork are sold which one would probably judge to be unrealistically high. Further, there are more than 36,000 hogs on the farm.

A row sum is simply a detailed breakdown of a constraint: each variable appearing in that constraint, its corresponding coefficient (a_{ij}) and the product $a_{ij}X_j$. The products are then summed, and subtracted from the right hand side and the slack variable formed. The use of row summing in our example begins with the pork sales constraint to see if 5.4 million lbs. is reasonable (Table 17.16.).

The pork constraint contains the variables sell pork and hog production. The sell pork variable uses one pound of pork per unit, while the hog production variable yields 150 pounds of pork per unit. The second column of Table 17.15 contains the optimal variable values. In the third column we write the product of the variable value and its a_{ij} . The products are summed to give total endogenous use which in this case equals zero. We then enter the right hand side and subtract it to determine the value of the slack variable. All these items in this case are zero. Given institutional knowledge, one would conclude the error has not yet been found as the 150 lbs. of pork per hog is reasonable, and all pork produced is sold. However, one would wonder if a production level of 36,001 hogs is reasonable. The next step is to examine the resources used by hog production. For illustrative purposes, we begin with the miscellaneous input supply-demand balance. The row sum for this constraint is shown in Table 17.17.

There are four entries in the constraint involving both basic and nonbasic variables. The row sum does not reveal anything terribly unrealistic except the large amount of activity from the hog production variable. The basic question is yet to be resolved.

We next investigate the corn supply-demand balance. The row sum computations for this constraint are shown in Table 17.18. In this case the constraint has a non-zero right hand side;

thus, the endogenous sum is 20 which equals the right hand side leaving the slack variable zero. We find the 36,001 hogs require 720,020 bushels of corn, and the reason they are able to obtain all this corn is because of the inaccurate yield on the corn production variable. The modeler would then correct the yield on the corn production variable.

The above example illustrates the principles behind using the allocation results to debug a model. One identifies a variable or slack with an unrealistically high solution value, and then row sums the constraints in which that variable is involved with to discover the problem. Row summing can be used to discover why particular variables have unrealistically large values by identifying incorrect coefficient values or coefficient placement errors. For example, suppose that the corn yield was inadvertently punched in the soybean row; then one might have discovered a solution in which soybeans are sold but no soybeans are produced. A row sum would quickly determine the source of the soybeans and indicate the error. Row summing can also be applied to discover the causes of large values for slack or surplus variables.

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Table 17.1. Priorities of Techniques to Use to Diagnose Improper Model Solution Outcomes

of Solution Outcome	Type	Structural Check ^a	Degen. Resol.	Scaling ^a	Artificial Variables	Upper Bounds	Budget	Row Sum
Solver Failure		1	3	2	5	4		
Unbounded Solution		1		3		2	4	
Infeasible Solutions		1		3	2		4	5
Unsat. Optimal Solutions		1					2	2

Notes: The entries in the table gives information on the order in which to try techniques with the technique numbered 1 being the item to try first.

^a This technique could be employed before any solving occurs. The technique also can be used when problems appear.

Table 17.2. Solution Properties of Various Model Formulations

Cases Where the Model Must have an Infeasible Solution

$b_i < 0$ and $a_{ij} \geq 0$ for all $j \in \Psi$ row i will not allow a feasible solution

$d_n < 0$ and $e_{nj} \geq 0$ for all $j \in \Psi$ row n will not allow a feasible solution

$d_n > 0$ and $e_{nj} \leq 0$ for all $j \in \Psi$ row n will not allow a feasible solution

$g_m > 0$ and $f_{mj} \leq 0$ for all $j \in \Psi$ row m will not allow a feasible solution

Cases where certain variables in the model must equal zero

$b_i = 0$ and $a_{ij} \geq 0$ for all $j \in \Psi$ all X_j 's with $a_{ij} \neq 0$ in row i will be zero

$d_n = 0$ and $e_{nj} \geq 0$ for all $j \in \Psi$ all X_j 's with $e_{nj} \neq 0$ in row n will be zero

$d_n = 0$ and $e_{nj} \leq 0$ for all $j \in \Psi$ all X_j 's with $e_{nj} \neq 0$ in row n will be zero

$g_m = 0$ and $f_{mj} \leq 0$ for all $j \in \Psi$ all X_j 's with $f_{mj} \neq 0$ in row m will be zero

Cases where certain constraints are obviously redundant

$b_i \geq 0$ and $a_{ij} \leq 0$ for all j means row i is redundant

$g_m \leq 0$ and $f_{mj} \geq 0$ for all j means row m is redundant

Cases where certain variables cause the model to be unbounded

$c_j > 0$ and $a_{ij} \leq 0$ or $e_{nj} = 0$ and $f_{mj} \geq 0$ for all $i, m,$ and n means variable j is unbounded

Cases where certain variables will be zero at optimality

$c_j < 0$ and $a_{ij} \geq 0$ or $e_{nj} = 0$ and $f_{mj} \leq 0$ for all $i, m,$ and n means variable j will always be zero

Table 17.3. Relationships Between Items Before and After Scaling

Item	Symbol Before Scaling	Symbol After Scaling	Unscaled Value in Terms of Scaled Value	Scaled Value in Terms of Unscaled Value
Variables	X_j	X_j'	$X_j = X_j' * (\text{COLSCAL}_j * \text{RHSSCAL})$	$X_j' = X_j / (\text{COLSCAL}_j * \text{RHSSCAL})$
Slacks	S_i	S_i'	$S_i = S_i' * (\text{ROWSCAL}_i * \text{RHSSCAL})$	$S_i' = S_i / (\text{ROWSCAL}_i * \text{RHSSCAL})$
Reduced Cost	$z_j - c_j$	$z_j' - c_j'$	$z_j - c_j = (z_j' - c_j') * (\text{OBJSCAL} / \text{COLSCAL}_j)$	$z_j' - c_j' = (z_j - c_j) / (\text{OBJSCAL} / \text{COLSCAL}_j)$
Shadow Price	U_i	U_i'	$U_i = U_i' * (\text{OBJSCAL} / \text{ROWSCAL}_i)$	$U_i' = U_i / (\text{OBJSCAL} / \text{ROWSCAL}_i)$
Obj. Func. Value	Z	Z'	$Z = Z' * \text{OBJSCAL} * \text{RHSSCAL}$	$Z' = Z / (\text{OBJSCAL} * \text{RHSSCAL})$

Table 17.4. Optimal Solution to Unscaled Nibble Production Problem

Obj = 3,600,000							
Variable	Units	Value	Reduced Cost	Equation	Unit	Slack	Shadow Price
X ₁	Lbs. of Nibbles	4000000	0	1	Lbs. of Nibbles	0	1
X ₂	Hrs. of Process 1	400	0	2	Lbs. of Gribbles	0	100
X ₃	Hrs. of Process 2	0	4800	3	# of Hibbles	0	6
X ₄	Sacks of Gribbles	40	0	4	Hrs of Labor	10000	0

Table 17.5. Optimal Solution to Nibble Production Problem After Row Scaling

Variable	Units	Value	Reduced Cost	Equation	Unit	Slack	Shadow Price
X ₁	Lbs. of Nibbles	4000000	0	1	1000's of Lbs. of Nibbles	0	1000
X ₂	Hrs. of Process 1	400	0	2	Lbs. of Gribbles	0	100
X ₃	Hrs. of Process 2	0	4800	3	100's of Hibbles	0	600
X ₄	Sacks of Gribbles	40	0	4	10's of Hrs of Labor	10000	0

Table 17.6. Optimal Solution to Nibble Production Problem After Row and Column Scaling

Variable	Units	Value	Reduced Cost	Equation	Unit	Slack	Shadow Price
X ₁	1000's of Lbs. of Nibbles	4000000	0	1	1000's of Lbs. of Nibbles	0	1000
X ₂	Hrs. of Process 1	400	0	2	Lbs. of Gribbles	0	100
X ₃	Hrs. of Process 2	0	4800	3	100's of Hibbles	0	600
X ₄	Sacks of Gribbles	40	0	4	10's of Hrs of Labor	10000	0

Table 17.7. Optimal Solution to Nibble Production Problem After Row, Column, Objective Function and RHS Scaling

Variable	Units	Value	Reduced Cost	Equation	Unit	Slack	Shadow Price
X ₁	100,000's Lbs. of Nibbles	40	0	1	100,000's of Lbs. of Nibbles	0	1
X ₂	100's of Hrs. of Process 1	4	0	2	100's Lbs. of Gribbles	0	0.1
X ₃	100's of Hrs. of Process 2	0	4.8	3	10,000's of Hibbles	0	0.6
X ₄	100's of Sacks of Gribbles	20	0	4	1000's of Hrs of Labor	100	0

Table 17.8. Solution to Infeasible Example with Artificial Present

Objective Function = -186935					
Variable	Value	Reduced Cost	Equation	Level	Shadow Price
X ₁	1.3	0	1	48.7	0
X ₂	0	1940	2	0	1990
A	18.7	0	3	0	-10,000

Table 17.9. Tableau of Budgeting Example								
Row	Buy Misc.	Sell Corn	Sell Soyb.	Sell Pork	Prod Corn	Prod Soyb.	Prod Hogs	RHS
Objective Func	-1	2.5	6	0.5	-75	-50		MAX
Land Available					1	1		≤ 600
Pork Balance				1			-1000	≤ 0
Soybean Bal			1			-50		≤ 0
Corn Balance		1			-120		20	≤ 0
Misc. Inp. Bal.	-1				125	50	20	≤ 0

Table 17.10. Optimal Solution to Budgeting Example

Variable	Value	Reduced Cost	Equation	Level	Shadow Price
Buy Misc. Input	147,000	0	Land Available	0	2680.00
Sell Corn	0	22.50	Pork Balance	0	0.5
Sell Soybeans	0	0	Soybean Balance	0	6.00
Sell Pork	3,600,000	0	Corn Balance	0	24.00
Produce Corn	600	0	Misc. Input Balance	0	1.00
Produce Soybeans	0	2,480.00			
Produce Hogs	3,600	0			

Table 17.11. Budget of Soybean Production Activity

Constraint	a_{ij}	Shadow Price (U_i)	Product ($U_i a_{ij}$)
Land Available	1	2680	2680
Soybean Balance	-50	6	-300
Misc. Input Balance	50	1	50
Indirect Cost Sum ($\sum U_i a_{ij}$)			2430
Less Objective Function (c_j)	-50		-(-50)
Red. Cost ($\sum U_i a_{ij} - c_j$)			2480(nonbasic)

Table 17.12. Corn Production Budget

Constraint	a_{ij}	Shadow Price (U_i)	Product ($U_i a_{ij}$)
Land Available	1	2680	2680
Corn Balance	-120	24	-2880
Misc. Input Balance	125	1	125
Indirect Cost Sum ($\sum U_i a_{ij}$)			-75
Less Objective Function (c_j)	-75		-(-75)
Reduced Cost($\sum U_i a_{ij} - c_j$)			0(basic)

Table 17.13. Hog Production Budget

Constraint	a_{ij}	Shadow Price (U_i)	Product ($U_i a_{ij}$)
Pork Balance	-1000	0.5	-500
Corn Balance	20	24	480
Misc. Input Balance	20	1	20
Indirect Cost Sum ($\sum U_i a_{ij}$)			0
Less Objective Function (c_j)	0		-(0)
Reduced Cost ($\sum U_i a_{ij} - c_j$)			0(basic)

Row	Buy Misc.	Sell Corn	Sell Soyb.	Sell Pork	Prod Corn	Prod Soyb.	Prod Hogs	RHS
Objective Func	-1	2.5	6	0.5	-75	-50		MAX
Land Available					1	1		≤ 600
Pork Balance				1			-150	≤ 0
Soybean Bal			1			-50		≤ 0
Corn Balance		1			-1200		20	≤ 20
Misc. Inp. Bal.	-1				125	50	20	≤ 0

Table 17.15. Optimal Solution to Row Summing Example

Variable	Value	Reduced Cost	Equation	Level	Shadow Price
Buy Misc. Input	795,020	0	Land Available	0	3,100
Sell Corn	0	0.25	Pork Balance	0	0.5
Sell Soybeans	0	0	Soybean Balance	0	6.00
Sell Pork	5,400,150	0	Corn Balance	0	2.75
Produce Corn	600	0	Misc. Input Bal	0	1.00
Produce Soybeans	0	2,480.00			
Produce Hogs	36,001	0			

Table 17.16. Row Sum of Pork Constraint

Variable	a_{ij}	Optimal Value (X_j^*)	Product ($a_{ij}X_j^*$)
Sell Pork	1	5,400,150	5,400,150
Produce Hogs	-150	36,001	-5,400,150
Endogenous Sum ($\sum a_{ij} X_j^*$)			0
Right Hand Side(b_i)	0		0
Slack ($b_i - \sum a_{ij} X_j^*$)			0

Table 17.17. Row Sum of Miscellaneous Input Constraint

Variable	a_{ij}	Optimal Value (X_j^*)	Product ($a_{ij}X_j^*$)
Buy Miscellaneous Inputs	-1	795,020	-795,020
Produce Corn	125	600	75,000
Produce Soybeans	50	0	0
Produce Hogs	20	36,001	720,020
Endogenous Sum ($\sum a_{ij} X_j^*$)			0
Right Hand Side(b_i)	0		0
Slack ($b_i - \sum a_{ij} X_j^*$)			0

Table 17.18. Row Sum of Corn Balance Constraint

Variable	a_{ij}	Optimal Value (X_j^*)	Product ($a_{ij}X_j^*$)
Sell Corn	1	0	0
Produce Corn	-1,200	600	-720,000
Produce Hogs	20	36,001	720,020
Endogenous Sum ($\sum a_{ij} X_j^*$)			20
Right Hand Side(b_i)	20		20
Slack ($b_i - \sum a_{ij} X_j^*$)			0